

A path finding algorithm.

## A path finding algorithm.

- Given a state space, such as a 2-dimensional map, find a path from point $A$ to point $B$ in that space, if such a path exists.
- If such a path exists, return a path within certain criteria - ie: shortest path, most straight path, avoiding certain areas, etc.


## The A* Algorithm

- Developed in 1968 for solving different kind of problems such as the '15-puzzle'.
- Searches for the least costly path from a starting state to a goal state by examining adjacent states of a particular state.


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## Starting State Adjacent States Goal State



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- The algorithm works by repeatedly examining the most promising unexplored adjacent state.
- A priority queue 'Open' contains all adjacent unexamined states, sorted in order of lowest cost.
- A list ' ' contains all examined states.
- Initially, the list is empty, while the Open list contains a single starting state.


## A Simple Example



## A Simple Example



Adjacent States (1 move from start)

## A Simple Example



## A Simple Example



## A Simple Example



## A Simple Example



## A Simple Example



|  | Closed |
| :---: | :---: |
| B1 (1) | A1 (0) |
| A2 (1) |  |
|  |  |
|  |  |
|  |  |

Adjacent States (2 moves from start)

## A Simple Example

| 1 | A B C |  |  |  | Closed |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | B1 (1) | A1 (0) |
|  | 0 | 1 | 2 | A2 (1) |  |
|  | 1 |  |  | C1 (2) |  |
|  | $\downarrow$ |  |  | B2 (2) |  |
|  | 2 |  |  | A3 (2) |  |

## A Simple Example



## A Simple Example



## A Simple Example



## A Simple Example



Adjacent States (3 move from start)

## A Simple Example

|  | $A \quad B$ | c |  | Closed |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  | C1 (2) | A1 (0) |
| 1 | $0 \rightarrow 1 \rightarrow 2$ |  | B2 (2) | B1 (1) |
| 2 | $1 \rightarrow 2$ | 3 | A3 (2) | A2 (1) |
|  | $\downarrow$ |  | C2 (3) |  |
|  | 2 | $G$ | B3 (3) |  |

## A Simple Example

| 1 | A B | c |  | Closed |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | A1 (0) |
|  | $0 \rightarrow 1 \rightarrow$ |  |  | B1 (1) |
| 2 | $1 \rightarrow$ | 3 |  | A2 (1) |
|  |  |  | C2 (3) |  |
| 3 |  | $G$ | $\rightarrow$ B3 (3) |  |

## A Simple Example

|  | $A \quad B$ | c |  | Closed |
| :---: | :---: | :---: | :---: | :---: |
|  |  |  |  | A1 (0) |
| 1 | $0 \rightarrow 1 \rightarrow 2$ |  |  | B1 (1) |
| 2 | $1 \rightarrow 2$ | 3 |  | A2 (1) |
| 3 | $\downarrow$ |  | C2 (3) | C1 (2) |
|  | 2 | G | B3 (3) | ... |

## A Simple Example



## A Simple Example



Adjacent State is Goal!

## A Simple Example



- It is now possible to construct a path back to the starting point.


## A Simple Example

- This example is very simple, and every state within the $3 \times 3$ grid was explored. However, many optimizations can be made to the $\mathrm{A}^{*}$ algorithm to increase efficiency.
- One such optimization is to add a Heuristic to the cost of each state evaluated. A good heuristic will increase efficiency - up to 100\% in best cases.


## Heuristic

## Bad Heuristic

 Good HeuristicClosed Search Space


## Summary

- The $A^{*}$ algorithm will always return the most efficient path, if one exists.
- If there is no path, however, then the $A^{*}$ algorithm becomes inefficient, as all state space will be explored.
- The $A^{*}$ algorithm can be extended to support multiple goals and multiple start locations, and is generally very adaptable to many different problem domains, ranging from music to computer graphics.

